

An Optimal Buffer Assignment Method for Streaming Applications on Hybrid Memory with Volatile and Non-volatile RAM

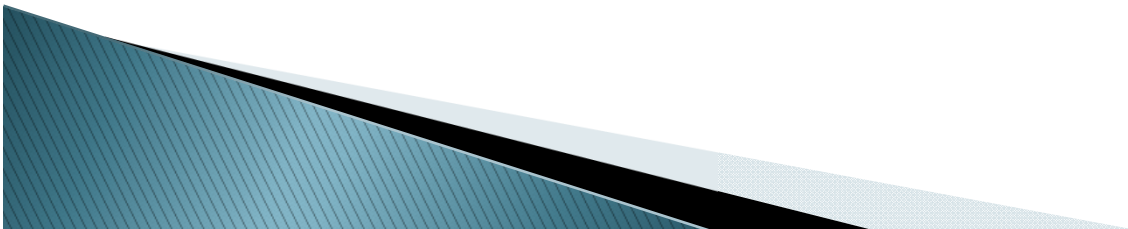
Hyunok Oh, Hanyang Univ.

2012.10.16

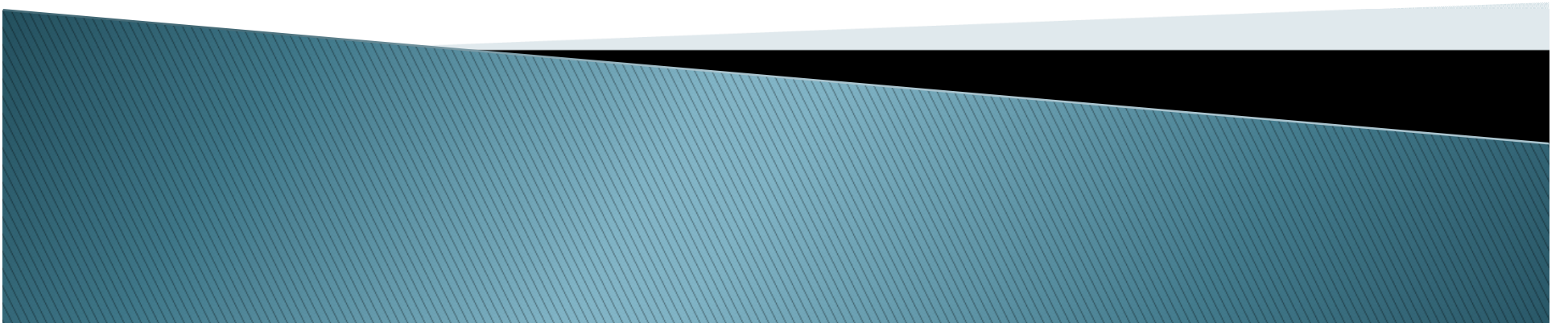


Outline

- ▶ Theoretical formulation
- ▶ Design tools



Theoretical Formulation



Outline

▶ Introduction

- Non-volatile Memory (NVM)
- Synchronous dataflow (SDF)

▶ Motivational Example

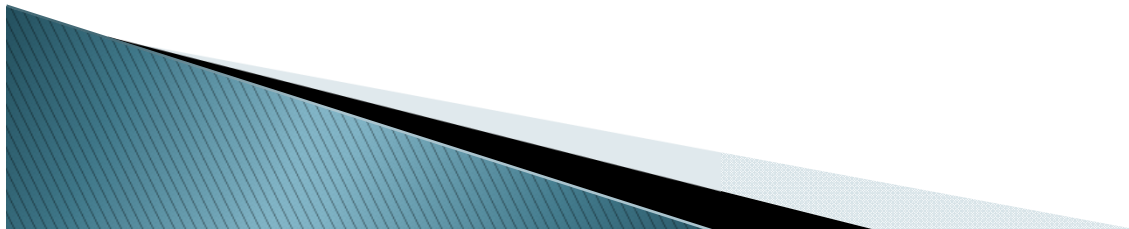
▶ Problem Definition

- P1. Maximization of PRAM life-time
- P2. Minimization of DRAM size
- P3. Minimization of energy consumption

▶ Answer Set Programming

▶ Experiment

▶ Conclusion

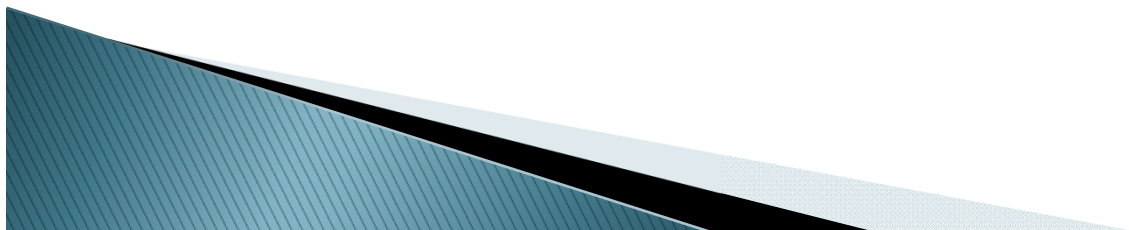


Introduction

▶ Non-Volatile Memory (NVM)

- Replace DRAM for main memory

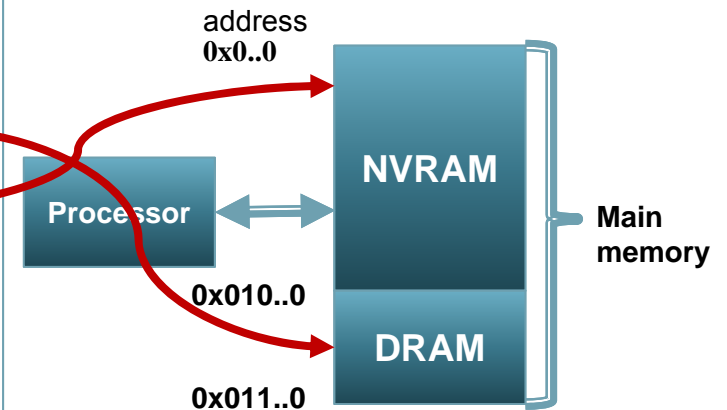
Type	Phase change RAM (PRAM) Spin-transfer torque magneto resistive RAM (STT-MRAM) Ferroelectric RAM(FRAM)
Pros	High density Low static energy consumption
Cons	Long write latency Short lifetime (PRAM : $10^5 \sim 10^8$ times)



Introduction

▶ Hybrid memory architecture – DRAM/NVRAM

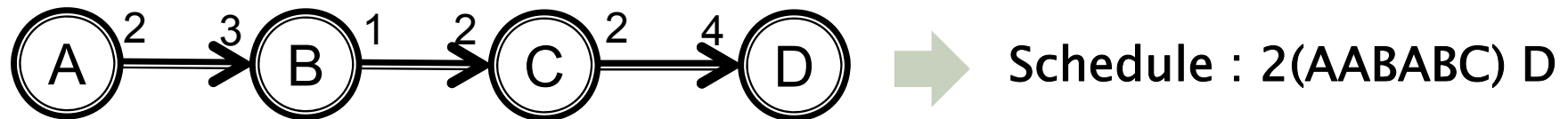
- Distributed data allocation
 - DRAM: frequently updated data
 - NVRAM: frequently read data
- Increase lifetime of NVRAM
- Reduce energy consumption



Introduction

▶ Synchronous dataflow (SDF)

- represents streaming applications like multimedia that require frequent memory access
- Node(Actor) - functional algorithm
- Edge - communication between two actors
- Producing / Consuming rate
 - the number of produced and consumed samples
- Rate is fixed



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▶ **Motivational Examples**

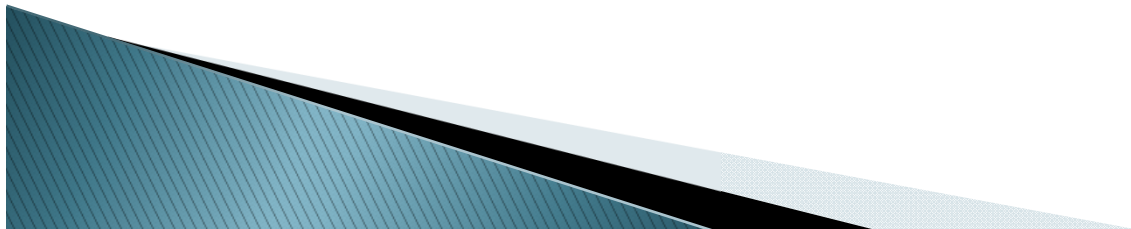
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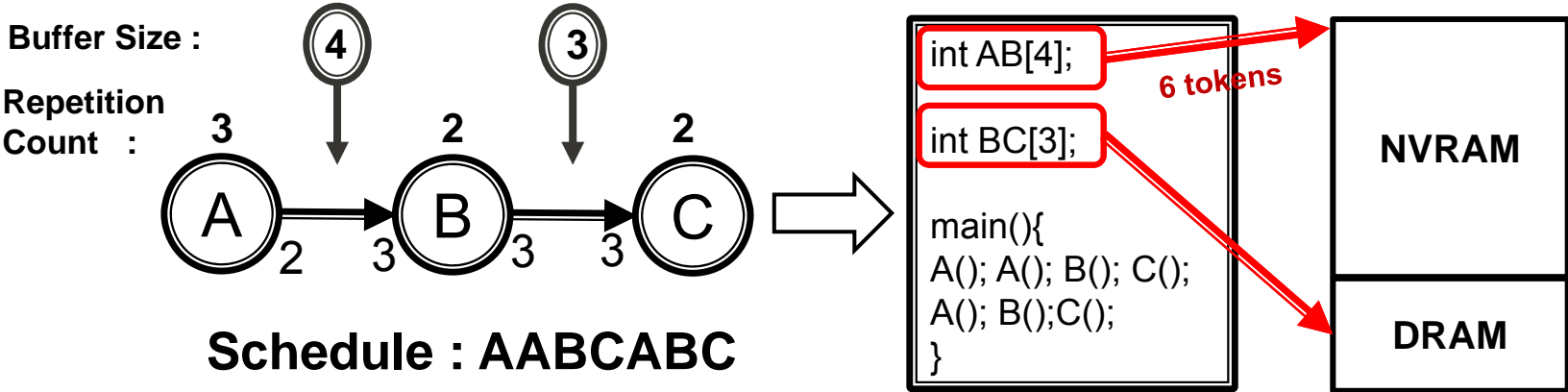
▶ Answer Set Programming

▶ Experiment

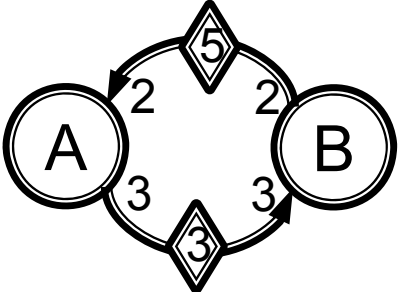
▶ Conclusion



Motivational Example



Motivational Example



Schedule : **AB**

Schedule : **BA**

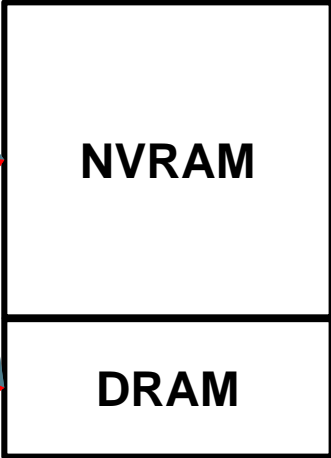
Minimizes the number of write to NVRAM

Minimal Buffer Schedule

```
int AB[6];  
int BA[5];  
main(){  
  A(); B();  
}
```

```
Int AB[7];  
Int BA[3];  
main(){  
  B(); A();  
}
```

Memory



2 tokens

3 tokens

Assumed Size : 6

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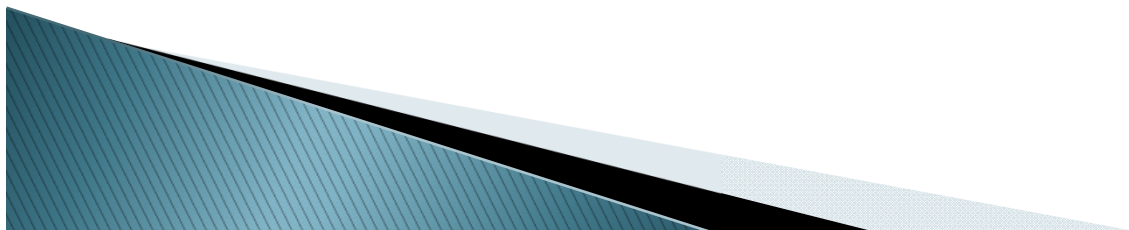
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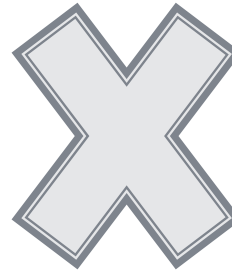


Problem Definition

(1) Maximize the lifetime of NVRAM memory

(2) Minimize the DRAM memory size

(3) Minimize the energy consumption



(a) For a given schedule

(b) By constructing an optimal schedule.

P1. Maximization of PRAM Life-time

► Formulation for given schedule

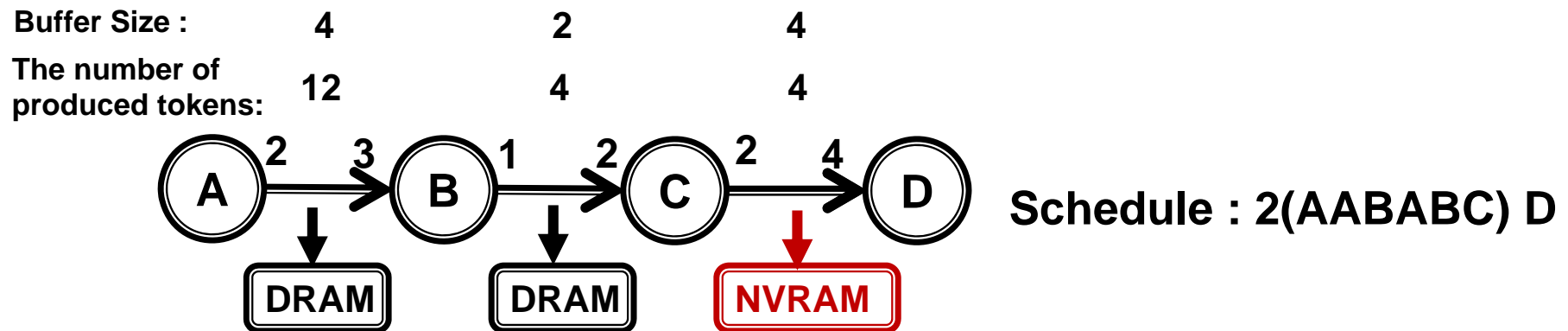
Minimize $\sum_{e_i \in E} p_i * s_i$ ← **The sum of the produced tokens assigned to NVRAM**

Subject to $\sum_{e_i \in E} d_i * b_i \leq S_{DRAM}$ ← **Total sum of the buffer size assigned to DRAM**

Constraint	Meaning
p_i	Binary variable to indicate whether the buffer on edge i is assigned to NVRAM or not ($p_i = 1 - d_i$)
s_i	The number of produced tokens on edge i per iteration
b_i	Binary variable to indicate whether the buffer on edge i is assigned to DRAM or not (It is 1 if assigned to DRAM)
d_i	The buffer size on edge i
E	A set of edges

P1. Maximization of PRAM Life-time

▶ Example



- Minimize $(12p_{AB}+4p_{BC}+4p_{CD})$
- Subject to $(4d_{AB}+2d_{BC}+4d_{CD}) \leq 6 (S_{DRAM})$
- ▶ $p_e + d_e$ (A buffer on each edge is assigned to either NVRAM or DRAM)

P1. Maximization of NVRAM Life-time

- ▶ Constructing an optimal schedule
 - Computes buffer size on each edge for each schedule to obtain an optimal schedule
 - Compared to previous example's optimal schedule
 - $2(3A2B)2CD$ - the number of produced tokens on NVRAM is **8**
 - Subject to $(6d_{AB}+2d_{BC}+4d_{CD}) \leq 6$
 - $6A\ 4B\ 2C\ D$ - the number of produced tokens on NVRAM is **16**
 - Subject to $(12d_{AB}+4d_{BC}+4d_{CD}) \leq 6$
 - **$2(AABABC)D$** - the number of produced tokens on NVRAM is **4**
 - Optimal schedule - $2(AABABC)D$
- ▶ It is applied to P2 and P3

P2.Minimization of DRAM Size

- ▶ Revise the endurance maximization problem

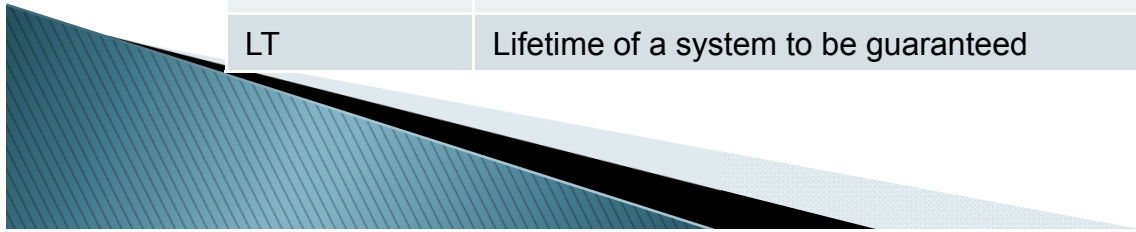
- Minimize $\sum_{e_i \in E} d_i * s_i$ The sum of the buffer size assigned to DRAM

- Subject to $\sum_{e_i \in E} p_i * s_i \leq \frac{P * R * S_{PRAM}}{LT}$

Total sum of the produced tokens assigned to NVRAM

A given minimum endurance of NVRAM

Constraint	Meaning
P	The period of an application
R	The endurance or the maximum reliable number of write operations per NVRAM cell.
LT	Lifetime of a system to be guaranteed



P3.Minimization of Energy Consumption

▶ Similarly to the minimization of DRAM Size

▶ Minimize $E_{NVRAM} \sum_{e_i \in E} p_i * s_i + E_{DRAM} \sum_{e_i \in E} d_i * s_i + S_{DRAM} * SE_{DRAM}$

Total energy consumption

Dynamic energy consumption on NVRAM

Dynamic energy consumption on DRAM

Static energy consumption on DRAM

▶ Subject to

$$\sum_{e_i \in E} p_i * s_i \leq \frac{P * R * S_{PRAM}}{LT}$$

NVRAM endurance

$$\sum_{e_i \in E} d_i * b_i \leq S_{DRAM}$$

DRAM Size

Constraint	Meaning
E_{NVRAM}, E_{DRAM}	Sum of energy consumptions of byte read and byte write per second on NVRAM, DRAM
SE_{DRAM}	Static energy consumption per byte on DRAM

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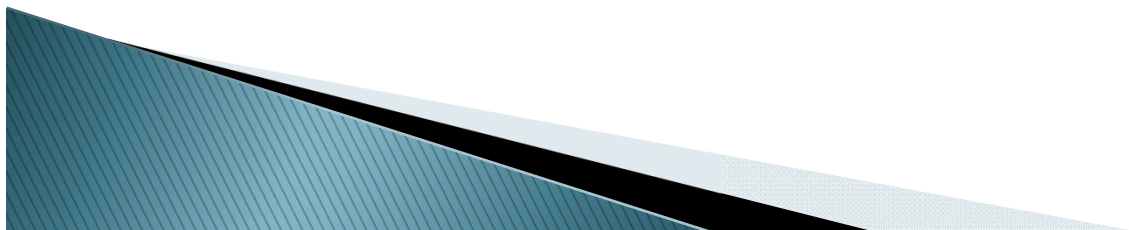
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▶ Answer Set Programming

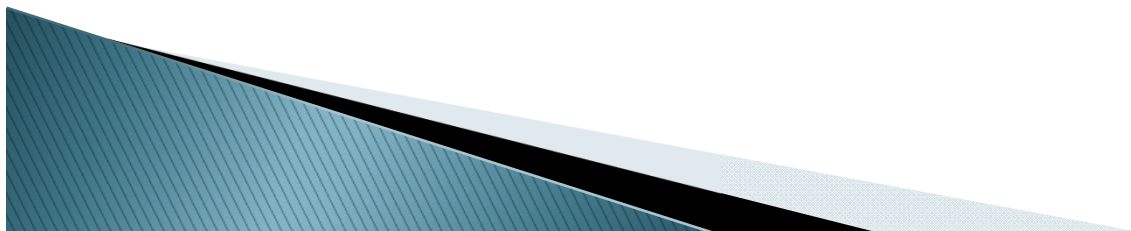
▶ Experiment

▶ Conclusion

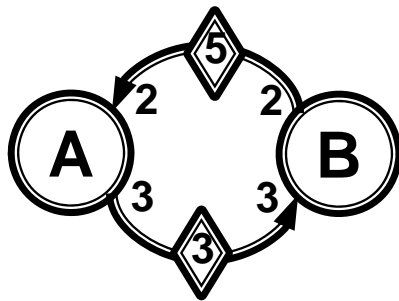


Answer Set Programming (ASP)

- ▶ **Declarative approach for NP problems**
- ▶ Problem - logic predicates “ **AND**”
- ▶ Solutions - answer sets
- ▶ Easy to understand the formulation
- ▶ Fast ASP solvers have been introduced



Problem Solving Flow



node(1..2).
 edge(1,1,2,3,3,3).edge(2,2,1,2,2,5).
 repetition(1,1).repetition(2,1).

SDF Graph



```
#begin_lua
function gcd(a,b)
if a=0 then return b
else return gcd(b%a, a)
end
end
#end_lua.
minimumBufSize(E,M) :- edge(E,A,B,P,C,I), M=P+C.@gcd(P,C).
bufSize(E,U) :- U = #max[sample(E,S)=S, minimumBufSize(E,M)=M], edge(E,_,_,_,_).
numSamples(E,S) :- S=P*X+I, edge(E,A,B,P,C,I), repetition(A,X).
0 { dram(E) : edge(E,_,_,_,_) }.
dramSum(E,S) :- dram(E), bufSize(E,S).
sample(E,I,0) :- edge(E,_,_,_,I).
sample(E,S+P,T) :- fire(A,T), edge(E,A,_,P,C,I), sample(E,S,T-1), time(T).
sample(E,S-C,T) :- fire(B,T), edge(E,_,B,P,C,I), sample(E,S,T-1), S>= C, time(T).
sample(E,S,T) :- sample(E,S,T-1), not fire(A,T), not fire(B,T), edge(E,A,B,_,_), time(T).
```

```
#sum[dramSum(E,S)=S : edge(E,_,_,_,_) : S=1..M : numSamples(E,M)] dramSize.
#maximize [dram(E)=N : numSamples(E,N)].
```

ASP Formulation



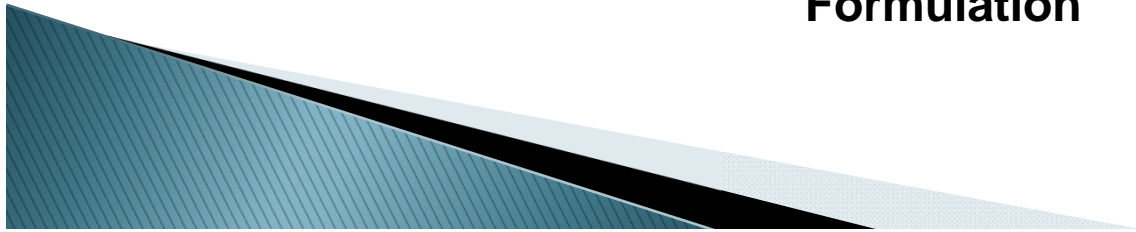
```
Answer: 1
bufSize(2,5),bufSize(1,6) numTotalSamplesOnPRAM(13)
Optimization: 13

Answer: 2
bufSize(2,5),bufSize(1,6)
dram(1) numTotalSamplesOnPRAM(7)
Optimization: 7

Answer: 3
bufSize(2,5),bufSize(1,6)
dram(2) numTotalSamplesOnPRAM(6)
Optimization: 6
OPTIMUM FOUND

Models : 1
Enumerated: 3
Optimum : yes
Optimization: 6
```

Result



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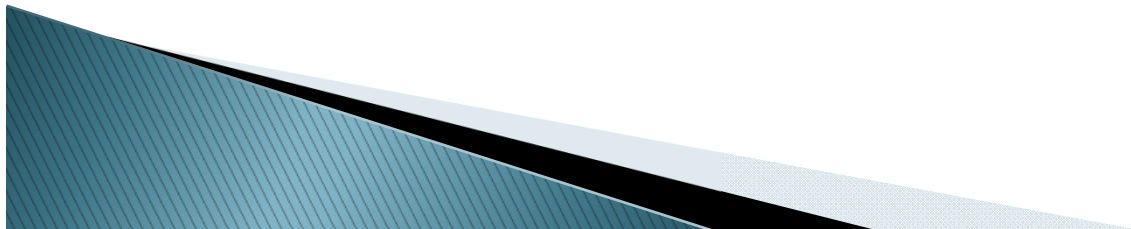
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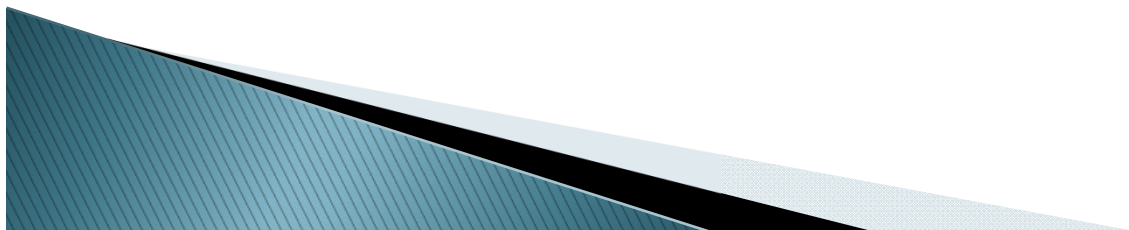


Experiments

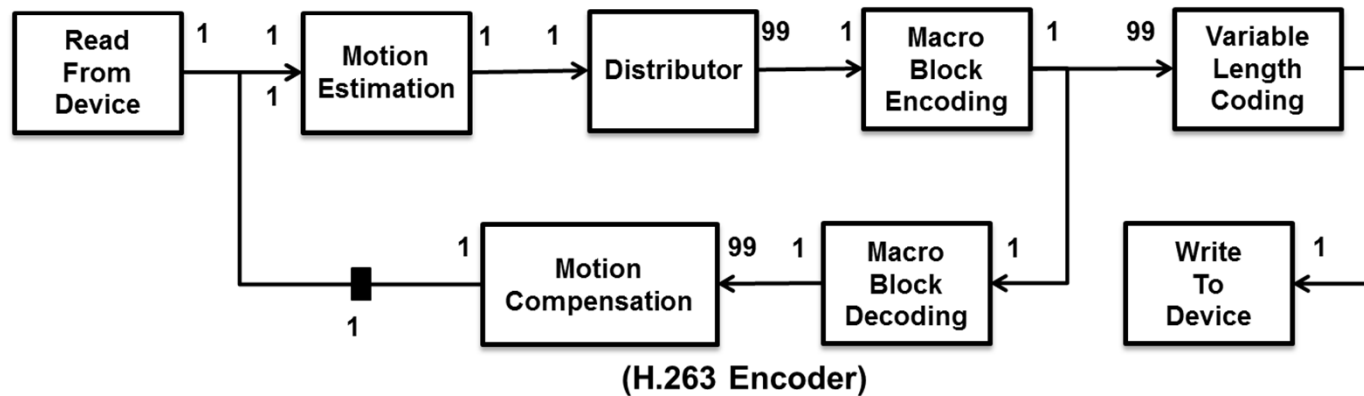
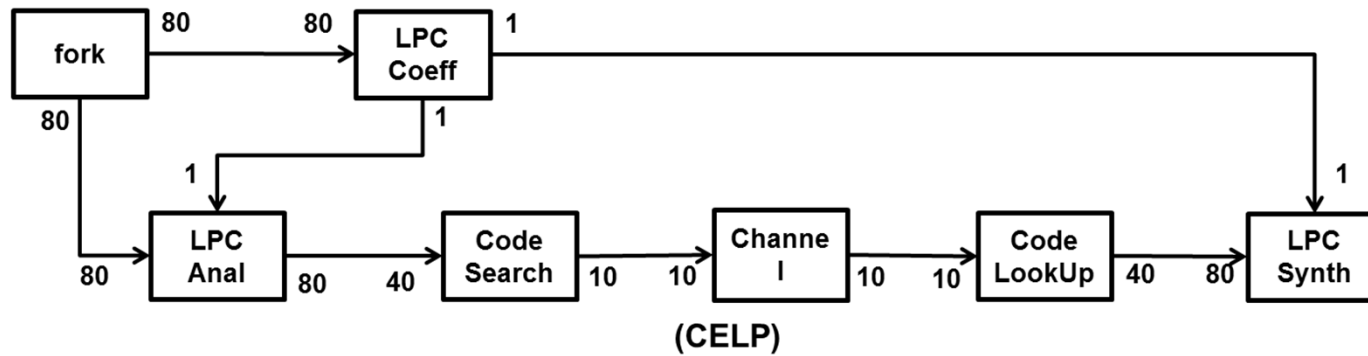
- ▶ Synthetic examples
- ▶ Real life applications
 - Part of CELP
 - H.263 encoder / decoder
 - MP3 decoder

CPU	Intel i7 2.7Ghz
RAM	6GB
OS	Ubuntu Linux
ASP Solver	Clingo 3.0

- ▶ 25 randomly generated examples
 - 3 to 6 actors
 - 1 to 3 repetition per actor



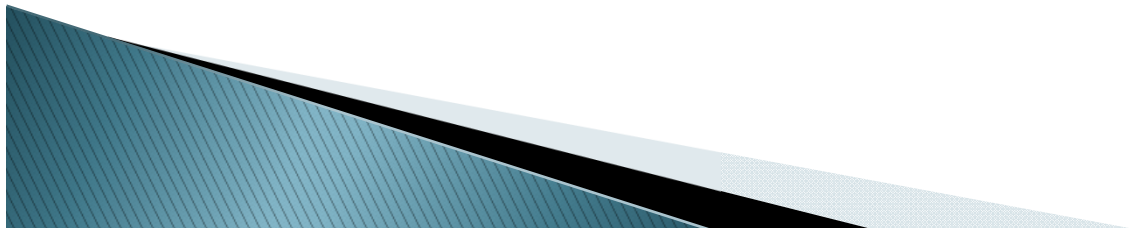
SDF Graphs



PRAM Lifetime Maximization for Real Life Applications

Graph	DRAM size	C1	C2	C3	C2 time	C3 time	$(C1-C2)/C2$
CELP	150	280	240	240	0.02	0.02	0.16
H.263 Encoder	150	523	396	396	0.42	30.82	0.32
H.263 Decoder	10	35	22	22	0.18	1.68	0.59
MP3	15	24	11	11	0.04	0.5	1.18

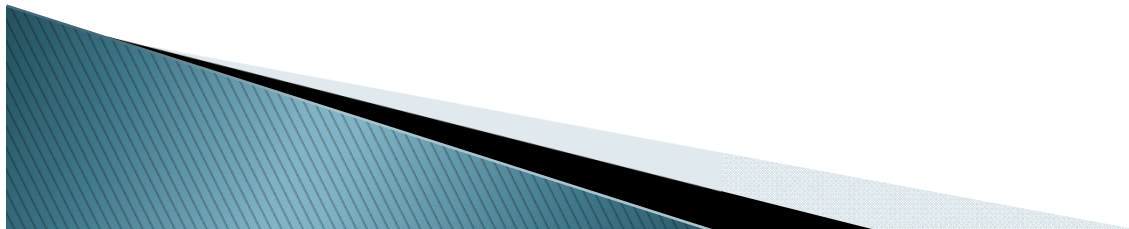
- ▶ The number of produced tokens
 - C1- Non-optimized approach with a given minimal buffer schedule
 - C2- Lifetime maximization with a given minimal buffer schedule
 - C3- Lifetime maximization with constructing an optimal schedule
- ▶ C2,C3 results
 - Acceptable to use existing scheduling algorithms to construct a schedule rather than constructing the optimal schedule
- ▶ Improvement of NVRAM lifetime by **63%** compared with a non-optimized algorithm



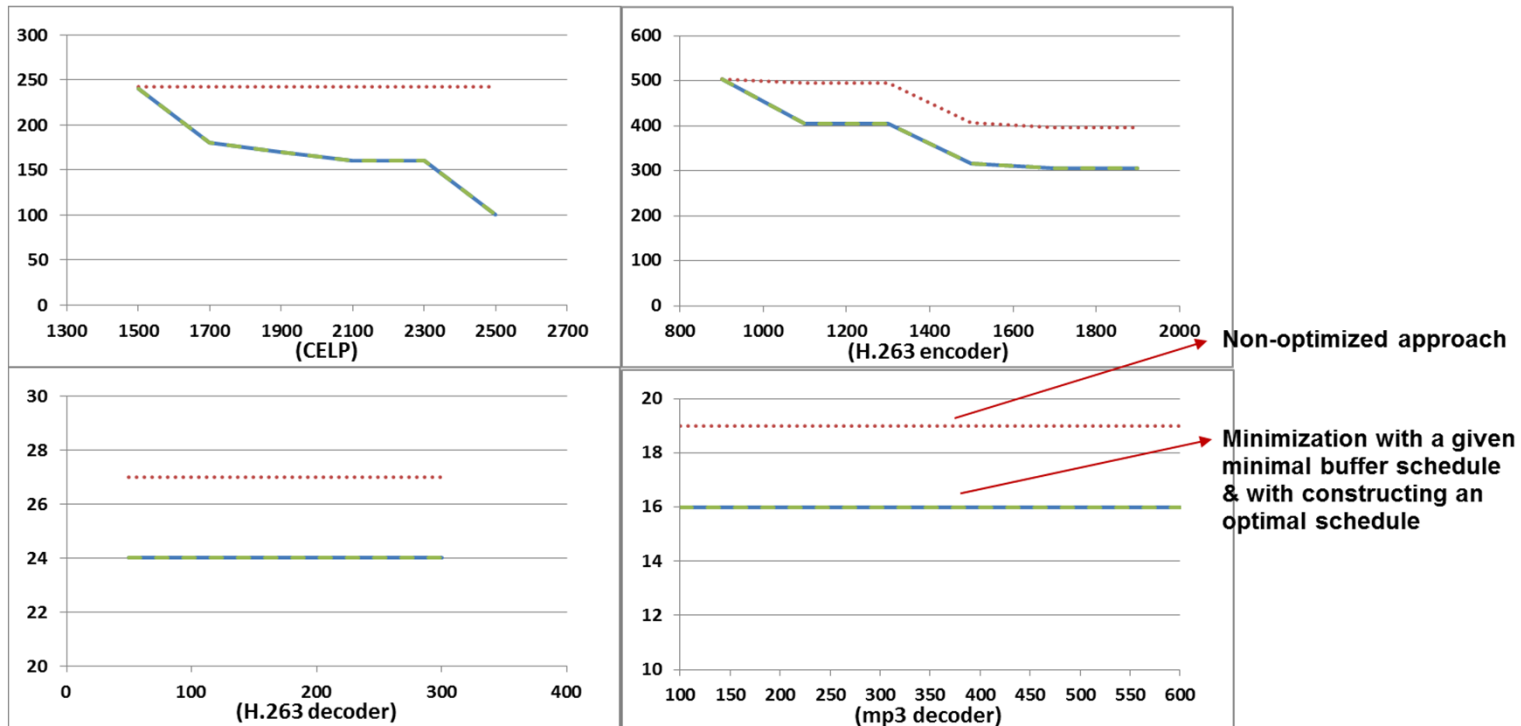
PRAM Lifetime Maximization for 25 Randomly Generated Examples

	$(C1-C2)/C2$	$(C2-C3)/C3$
Max. Lifetime	48.7%	0%
Min. DRAM Size	100.0%	0%

- ▶ Average improvement of NVRAM lifetime and the DRAM size with average computation times
- ▶ Similarly to the real-life applications, the optimizations with a given schedule and constructing an optimal schedule produce the same optimal results.

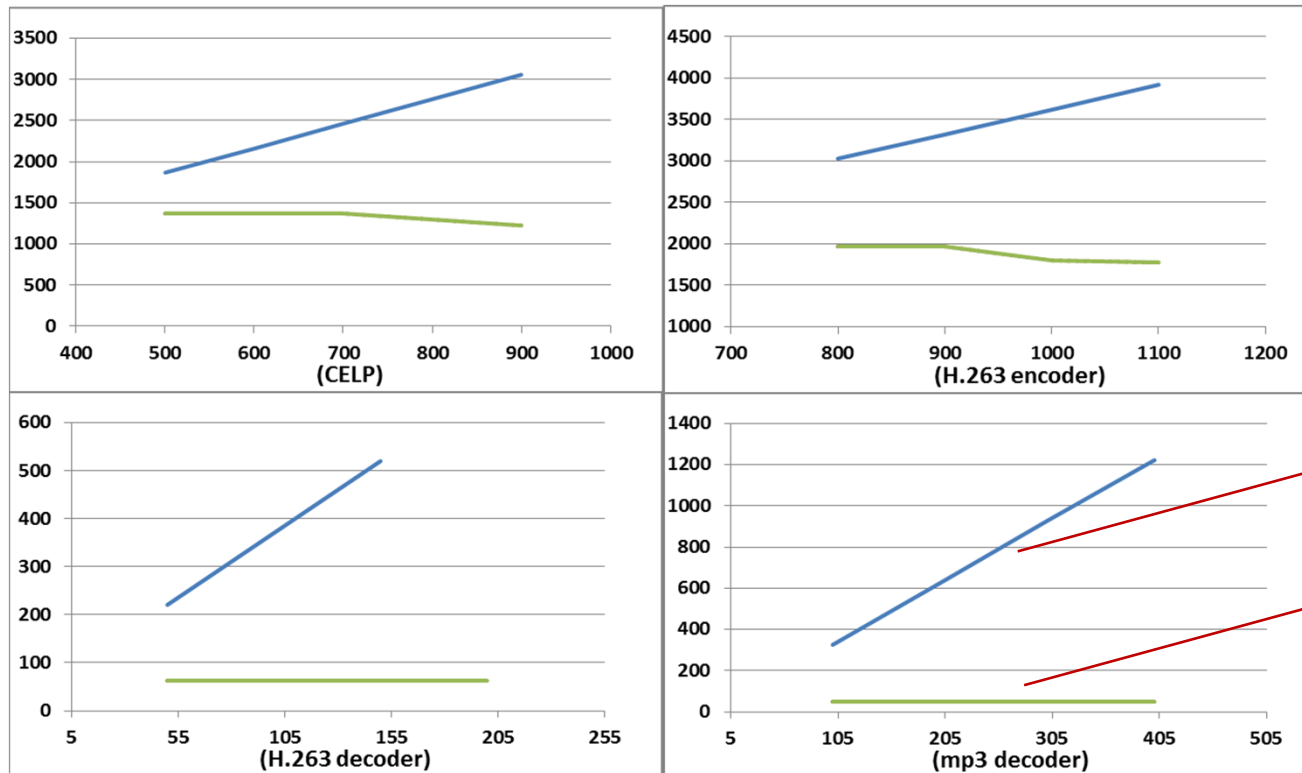


DRAM Size Minimization



- ▶ Minimization for a given NVRAM size with satisfying a given NVRAM lifetime constraint
- ▶ X axis - a given NVRAM size
- ▶ Y axis - the minimum DRAM size
- ▶ Minimization of DRAM size by 48.8% compared with a non-optimized algorithm

Energy Minimization



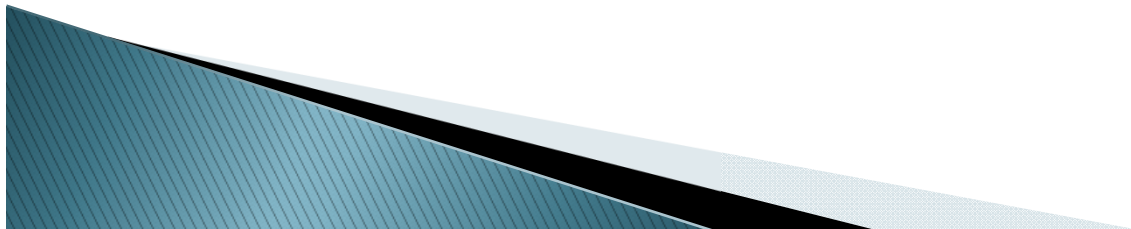
Non-optimized approach

Minimization with a given minimal buffer schedule & with constructing an optimal schedule

- ▶ Minimization of the energy consumption by varying the NVRAM size
- ▶ X axis : a given NVRAM size and
- ▶ Y axis : the minimal DRAM size
- ▶ Energy consumption saving by 165%

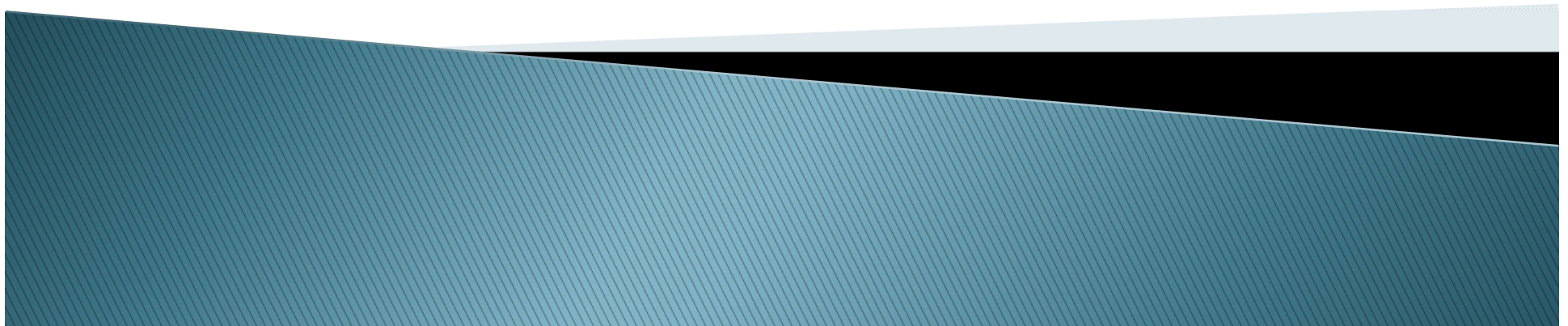
Summary

- ▶ Optimal buffer assignment algorithm for hybrid memory
- ▶ NVRAM lifetime maximization (63%)
- ▶ DRAM size minimization (48.8%)
- ▶ Energy consumption saving (165%)



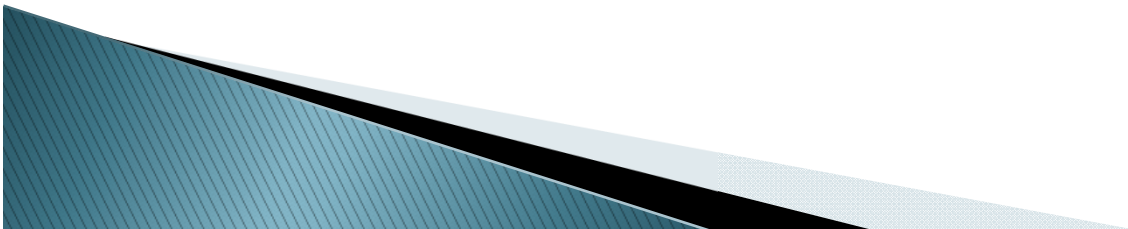
Design Tools :

Dataflow in Answer Set Programming (DASP)



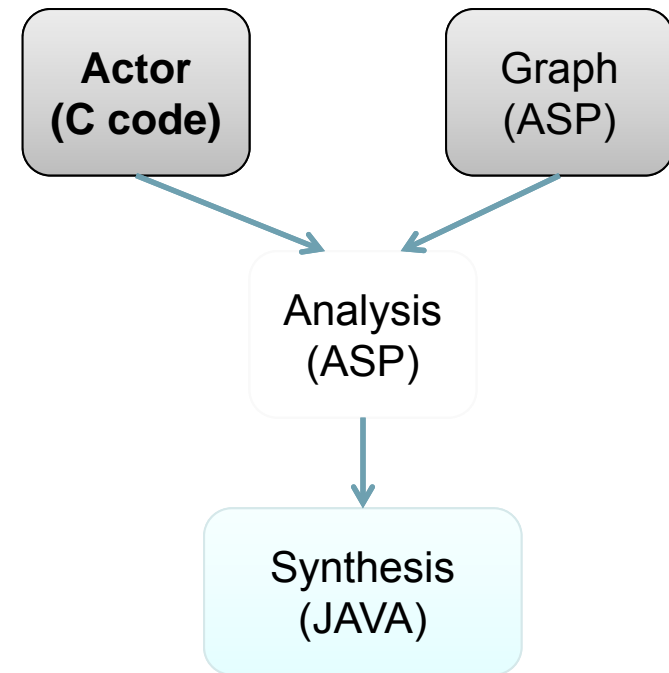
Goals of DASP

- ▶ Easy to use
 - No learning (C code based)
- ▶ Easy to install
 - Small size
- ▶ Easy to debug
 - C based
 - No code translation
- ▶ Easy to manage
 - Textual representation



DASP : Dataflow in Answer Set Programming

- ▶ Specification
 - Actor code
 - C code with macro keywords
 - Graph representation
 - Text (ASP)
- ▶ Mapping/Scheduling
 - Answer Set Programming
- ▶ Code synthesis
 - Java



DASP : Dataflow in Answer Set Programming

```
#include <stdio.h>

#define WIDTH 176
#define HEIGHT 144
```

```
typedef struct {
    char pixel[WIDTH*HEIGHT];
} Image;
```

```
INFO_BEGIN
    OUT_PORT(output,Image,2)
INFO_END
```

```
VARs_BEGIN
    int current;
VARs_END
```

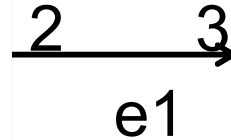
```
INIT_BEGIN
    VAR(current) = 0;
INIT_END
```

```
GO_BEGIN
    PORTn(output,0).pixel[0] = VAR(current);
    VAR(current)++;
    PORTn(output,1).pixel[0] = VAR(current);
    VAR(current)++;
```

```
    printf("send %d %d\n",PORTn(output,0).pixel[0],
PORTn(output,1).pixel[0]);
GO_END
```

```
WRAPUP_BEGIN
WRAPUP_END
```

ActorA.c



```
#include <stdio.h>

#define WIDTH 176
#define HEIGHT 144
```

```
typedef struct {
    char pixel[WIDTH*HEIGHT];
} Image;
```

```
INFO_BEGIN
    IN_PORT(input,Image,3)
INFO_END
```

```
VARs_BEGIN
VARs_END
```

```
INIT_BEGIN
INIT_END
```

```
GO_BEGIN
    printf("receive %d %d %d\n", PORTn(input,0).pixel[0],
PORTn(input,1).pixel[0], PORTn(input,2).pixel[0]);
GO_END
```

```
WRAPUP_BEGIN
WRAPUP_END
```

ActorB.c

```
schematic(prodcons).
actor(actorA). actor(actorB).
instance(a,actorA). instance(b,actorB).
connect(e1,a,output,b,input).
```

Graph

Analysis and Scheduling

- ▶ Answer Set Programming
- ▶ CreateGraph
- ▶ Repetition
- ▶ Schedule

Answer: 1

```
port_info(a,output,"Image",25344,output,2)
port_info(b,input,"Image",25344,input,3) schematic(prodcons) actor(actorA)
actor(actorB) instance(a,actorA) instance(b,actorB)
connect(e1,a,output,b,input) time(1) node(b) node(a) edge(e1,a,b,2,3,0)
numEdges(1) sdfedgegcd(e1,1) possibleRepetition(a,3,0)
possibleRepetition(b,2,0) possibleRepetition(a,3,1) possibleRepetition(b,2,1)
repetition(a,3) repetition(b,2) totalInvocations(5) time(2) time(3) time(4)
time(5) sample(e1,0,0) fire(b,5) fire(b,4) fire(a,3) fire(a,2) fire(a,1)
sample(e1,2,1) sample(e1,4,2) sample(e1,6,3) sample(e1,3,4)
sample(e1,0,5) notfireable(b,2) notfireable(b,1) buffer_size(e1,6)
index_management(b,input,linear) index_management(a,output,linear)
top_schedule(a,1) top_schedule(a,2) top_schedule(a,3) top_schedule(b,4)
top_schedule(b,5)
```

Code Generation

a.h

```
#define a_output_port(n) e1[n+e1_output_index]
#define DECLARE_BUFFERS Image e1[6];

#define DECLARE_INDEXES int e1_output_index;
#define INIT_INDEXES e1_output_index=0;
#define UPDATE_INDEXES e1_output_index=
(e1_output_index+2)%6;
```

b.h

```
#define b_input_port(n) e1[n+e1_input_index]
#define DECLARE_BUFFERS extern Image e1[6];

#define DECLARE_INDEXES int e1_input_index;
#define INIT_INDEXES e1_input_index=0;
#define UPDATE_INDEXES
e1_input_index=(e1_input_index+3)%6;
```

```
#include <stdio.h>
#define WIDTH 176
#define HEIGHT 144

typedef struct {
    char pixel[WIDTH*HEIGHT];
} Image;

INFO_BEGIN
OUT_PORT(output,Image,2)
INFO_END

VARS_BEGIN
int current;
VARS_END

INIT_BEGIN
VAR(current) = 0;
INIT_END

GO_BEGIN
PORTn(output,0).pixel[0] = VAR(current);
VAR(current)++;
PORTn(output,1).pixel[0] = VAR(current);
VAR(current)++;

printf("send %d %d\n",PORTn(output,0).pixel[0], PORTn(output,1).pi
GO_END

WRAPUP_BEGIN
WRAPUP_END
```

ActorA.c

```
void a_init(); void a_go(); void a_wrapup();
void b_init(); void b_go(); void b_wrapup();

int main()
{
    int prodcons_count0;
    a_init(); b_init();

    for(prodcons_count0=0;prodcons_count0<3;p
rodcons_count0++) {
        a_go();
        a_go();
        a_go();
        b_go();
        b_go();
    }
    a_wrapup(); b_wrapup();
}
```

main.c

PORTn(input,0).pixel[0], PORTn(input,1).pixel[0],

ActorB.c

Result

```
[hoh-ui-MacBook-Air:trunk/examples/prodcons] hoh% ./run_script
rm -f port_token_type.lp
java -classpath /Users/hoh/Documents/Research/devel/dasp/trunk/synthesis TaskInformation --type a actorA.c >> port_token_type.lp
java -classpath /Users/hoh/Documents/Research/devel/dasp/trunk/synthesis TaskInformation --type b actorB.c >> port_token_type.lp
java -classpath /Users/hoh/Documents/Research/devel/dasp/trunk/synthesis TaskInformation --buffer a port_token_type.lp graph.lp > a.h
java -classpath /Users/hoh/Documents/Research/devel/dasp/trunk/synthesis TaskInformation --buffer b port_token_type.lp graph.lp > b.h
java -classpath /Users/hoh/Documents/Research/devel/dasp/trunk/synthesis TaskInformation --info graph.lp > genInfo.c
gcc -g -c genInfo.c
gcc -g -o a_info.o -c actorA.c -DTASK_NAME=a -DTASK_INFO_GENERATION -I../include -include task.h -include a.h
gcc -g -o b_info.o -c actorB.c -DTASK_NAME=b -DTASK_INFO_GENERATION -I../include -include task.h -include b.h
gcc -o prodcons_info genInfo.o a_info.o b_info.o
% warning: iport/5 is never defined
% warning: initialDelay/2 is never defined
% warning: -index_management/3 is never defined
java -classpath /Users/hoh/Documents/Research/devel/dasp/trunk/synthesis TaskInformation --buffer_schedule a port_token_type.lp
schedule.result > a.h
java -classpath /Users/hoh/Documents/Research/devel/dasp/trunk/synthesis TaskInformation --buffer_schedule b port_token_type.lp
schedule.result > b.h
gcc -g -c prodcons.c
gcc -g -o a.o -c actorA.c -DTASK_NAME=a -I../include -include task.h -include a.h
gcc -g -o b.o -c actorB.c -DTASK_NAME=b -I../include -include task.h -include b.h
gcc -o prodcons prodcons.o a.o b.o
send 0 1
send 2 3
send 4 5
receive 0 1 2
receive 3 4 5
send 6 7
send 8 9
send 10 11
receive 6 7 8
receive 9 10 11
send 12 13
send 14 15
send 16 17
receive 12 13 14
receive 15 16 17
```

Support of Intermediate Port

```
INFO_BEGIN
  OUT_PORT(output,int,1)
  IN_PORT(input,int,1)
  IPORT(output,default,0)
  IPORT(input,default,1)
INFO_END

VARS_BEGIN
  int current;
VARS_END

INIT_BEGIN
  VAR(current) = 0;
INIT_END

GO_BEGIN
  int i;
  PORT(output) = VAR(current);
  for(i=0; i<3; i++) {
    SEND(default,0)
    RECEIVE(default,1)
    PORT(output) = PORT(input);
  }
  VAR(current)+=10;
GO_END

WRAPUP_BEGIN
WRAPUP_END
```

```
#include <stdio.h>

INFO_BEGIN
  IN_PORT(input,int,1)
  OUT_PORT(output,int,1)
INFO_END

INIT_BEGIN
INIT_END

GO_BEGIN
  PORT(output) = PORT(input)*2;
  printf("actor B : %d\n",PORT(input));
GO_END

WRAPUP_BEGIN
WRAPUP_END
```

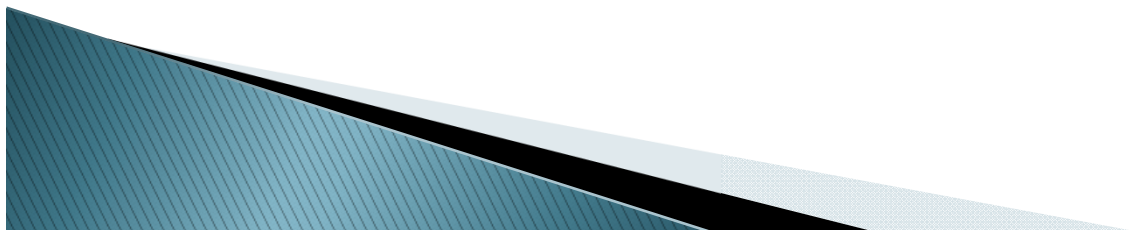


Running Result

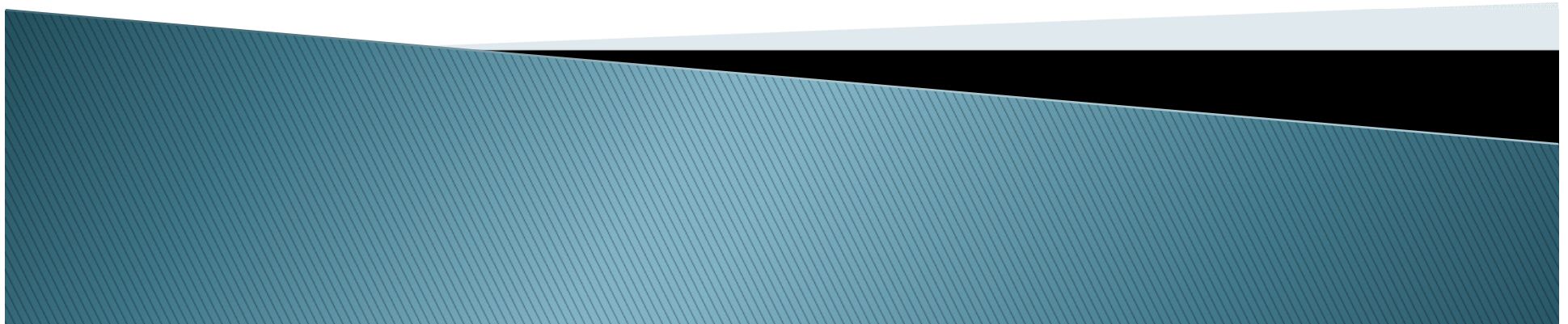
```
gcc -g -o closestCentroid_info.o -c closestCentroid.c -DTASK_NAME=closestCentroid -DTASK_INFO_GENERATION -I../include -include task.h -include closestCentroid.h
gcc -g -o updateCentroid_info.o -c updateCentroid.c -DTASK_NAME=updateCentroid -DTASK_INFO_GENERATION -I../include -include task.h -include updateCentroid.h
gcc -o kmeans_info genInfo.o readFile_info.o fork_point_info.o closestCentroid_info.o updateCentroid_info.o fork_point2_info.o print_info.o
% warning: iport/5 is never defined
% warning: -index_management/3 is never defined
java -classpath /Users/hoh/Documents/Research/devel/dasp/trunk/synthesis TaskInformation --buffer_schedule readFile port_token_type.lp schedule.result > readFile.h
java -classpath /Users/hoh/Documents/Research/devel/dasp/trunk/synthesis TaskInformation --buffer_schedule fork_point port_token_type.lp schedule.result > fork_point.h
java -classpath /Users/hoh/Documents/Research/devel/dasp/trunk/synthesis TaskInformation --buffer_schedule closestCentroid port_token_type.lp schedule.result > closestCentroid.h
java -classpath /Users/hoh/Documents/Research/devel/dasp/trunk/synthesis TaskInformation --buffer_schedule updateCentroid port_token_type.lp schedule.result > updateCentroid.h
java -classpath /Users/hoh/Documents/Research/devel/dasp/trunk/synthesis TaskInformation --buffer_schedule fork_point2 port_token_type.lp schedule.result > fork_point2.h
java -classpath /Users/hoh/Documents/Research/devel/dasp/trunk/synthesis TaskInformation --buffer_schedule print port_token_type.lp schedule.result > print.h
gcc -g -c kmeans.c
gcc -g -o closestCentroid.o -c closestCentroid.c -DTASK_NAME=closestCentroid -I../include -include task.h -include closestCentroid.h
gcc -g -o updateCentroid.o -c updateCentroid.c -DTASK_NAME=updateCentroid -I../include -include task.h -include updateCentroid.h
gcc -o kmeans kmeans.o readFile.o fork_point.o closestCentroid.o updateCentroid.o fork_point2.o print.o
centroid[0].p[0]=7.000000
centroid[0].p[1]=9.000000
centroid[1].p[0]=3.000000
centroid[1].p[1]=8.000000
centroid[2].p[0]=0.000000
centroid[2].p[1]=2.000000
centroid[3].p[0]=4.000000
centroid[3].p[1]=8.000000
output point 0.000000 0.000000
output point 1.000000 1.000000
output point 2.000000 2.000000
output point 3.000000 3.000000
output point 4.000000 4.000000
output point 5.000000 5.000000
output point 6.000000 6.000000
output point 7.000000 7.000000
output point 8.000000 8.000000
output point 9.000000 9.000000
read point 0.000000 0.000000
minIndex : 2, minDistance 2.000000
read point 1.000000 1.000000
minIndex : 2, minDistance 1.414214
read point 2.000000 2.000000
minIndex : 2, minDistance 2.000000
read point 3.000000 3.000000
minIndex : 2, minDistance 3.162278
read point 4.000000 4.000000
minIndex : 3, minDistance 4.000000
read point 5.000000 5.000000
minIndex : 3, minDistance 3.162278
read point 6.000000 6.000000
minIndex : 3, minDistance 2.828427
read point 7.000000 7.000000
minIndex : 0, minDistance 2.000000
read point 8.000000 8.000000
minIndex : 0, minDistance 1.414214
read point 9.000000 9.000000
minIndex : 0, minDistance 2.000000
Centroids
8.000000 8.000000
Centroids
0.000000 0.000000
Centroids
1.500000 1.500000
Centroids
5.000000 5.000000
```

Project : DASP

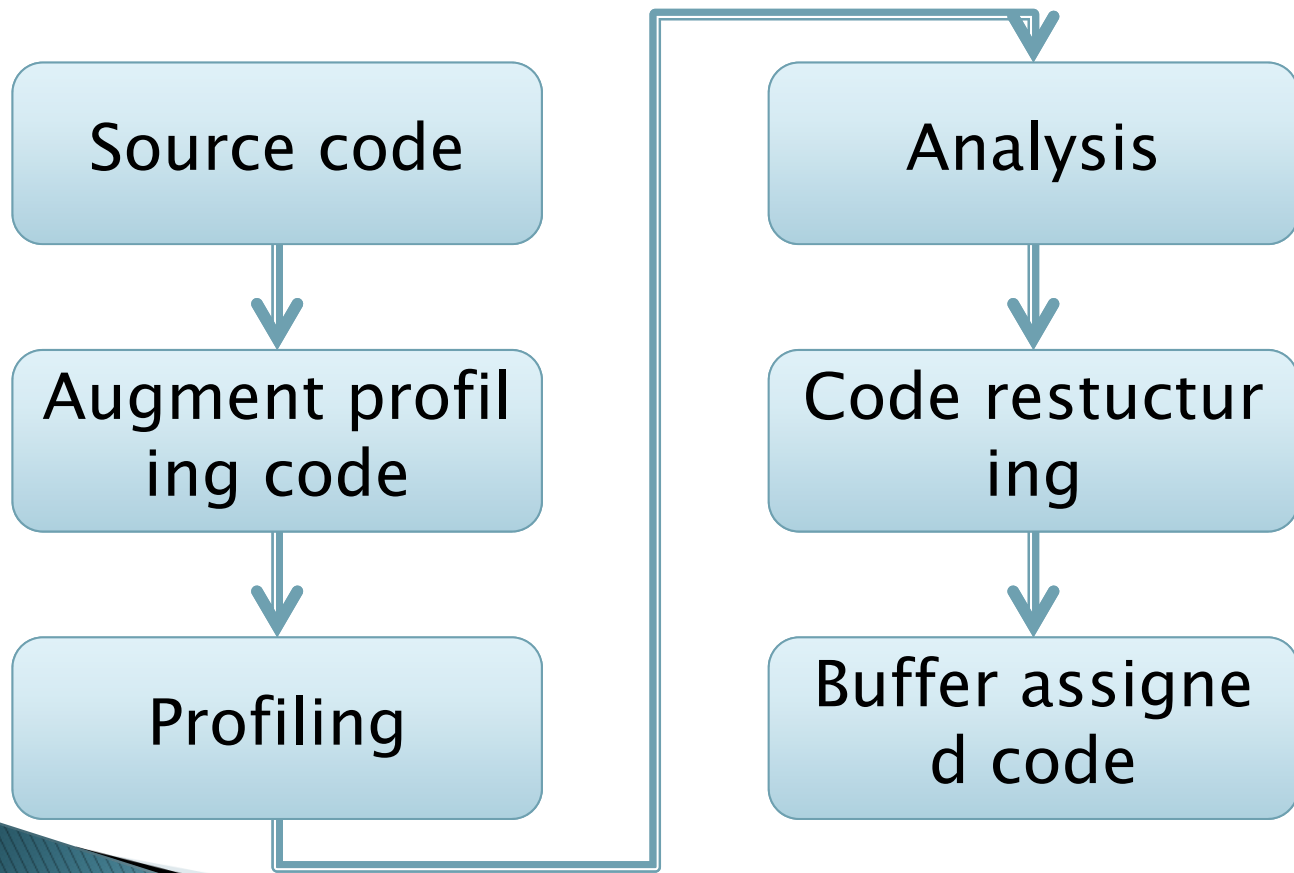
- ▶ <http://dev.naver.com/projects/dasp>
- ▶ On-going works
 - Support multi-core
 - Support distributed computing (MapReduce ...)
 - Implement multimedia applications
 - H.263 encoder/decoder, H.264/AVC encoder/decoder, HeVC encoder/decoder
 - Implement data mining applications
 - Clustering algorithm
 - Integrate all the developed algorithms
 - Buffer sharing, buffer index management, scheduling, hybrid memory



Design Tools : Compiler Approach



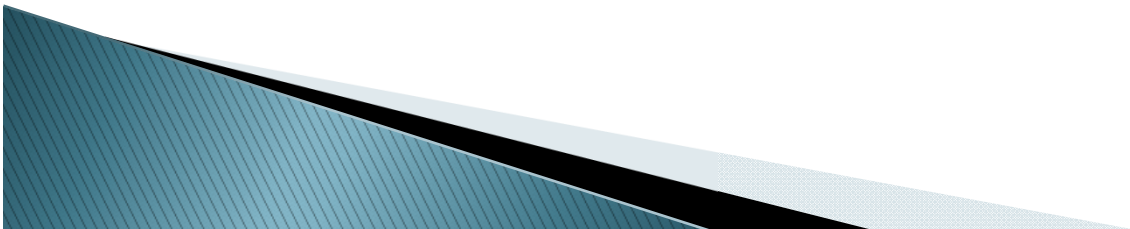
Buffer Assignment onto Hybrid Memory for FFMpeg



Read intensive :
NVRAM
Write intensive :
DRAM

Profiling

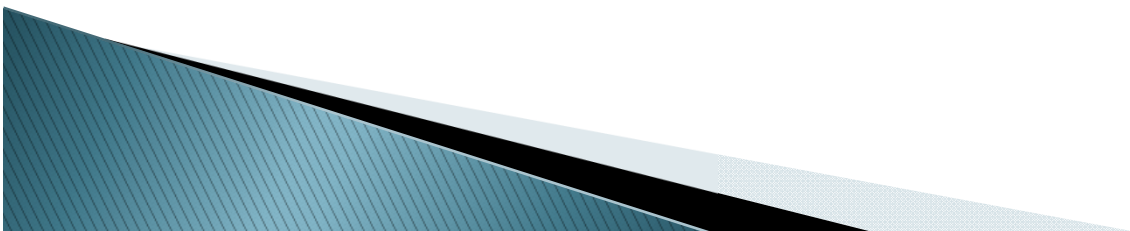
- ▶ MpegEncContext memory size : 10,468 Bytes
 - Total read Bytes : 93,418,145
 - Total write Bytes: 32,423,460



<u>MpegEncContext struct Members</u>	<u>Type</u>	<u>Read Count</u>	<u>Write Count</u>
->v_edge_pos	int	447201	
->current_picture.f.linesize[o]	int	894402	
->current_picture.f.linesize[o]	int	547200	
->mb_x	int	447201	
->mb_x	int	447201	
->mb_x	int	547200	
->mb_y		447201	
->mb_y	int	447201	
->mb_y	int	547200	
->mb_stride	int	447201	
->mb_stride	int	547200	
->mb_skipped	int		376931
->mb_skipped	int		376931
->interlaced_dct	int	547200	
->dest[0]	uint8_t	547200	
->mv_dir	int		376931
->mv_type	int		547200
->mv_type	int		447201
->block	short	1093792	
->linesize	int		
->y_dc_scale	int	399996	
->c_dc_scale	int	199998	
->current_picture.f.qscale_table	int8_t	599994	
->dsp.idct_permutation[]	uint8_t	8399916	
->h263_pred	int	547200	
->block_last_index[]	int		2261586
->current_picture.f.mb_type[]	uint32_t		547200
->mcsel	int		376931
->mv[][][]	int		517471
->mb_intra	int		119869
->ac_pred	int	99999	99999
->use_intra_dc_vlc	int		99999

Splitting the MpegEncContext struct

- ▶ MpegEncContext is most used struct in the FFmpeg application.
- ▶ Manually placing variable access counters.
- ▶ Move read intensive variable to second new struct
- ▶ Second struct is pre-allocated in memory in the global segment.
- ▶ Pre-allocated struct pointer is assigned to the MpegEncContext at the point of initialization.



MpegEncContext with pointer to second struct

```
typedef struct MpegEncContext {  
    .  
    .  
    MpegEncContext_Persistent_ * Persistent_ ;  
}  
MpegEncContext;
```

Second struct with read intensive members

```
typedef struct MpegEncContext_Persistent_{  
    int h_edge_pos, v_edge_pos  
    int mb_x, mb_y;  
    int mb_stride;  
    DCTELEM (*block)[64];  
    int y_dc_scale, c_dc_scale;  
    int h263_pred;  
    int ac_pred;  
}MpegEncContext_Persistent_;
```

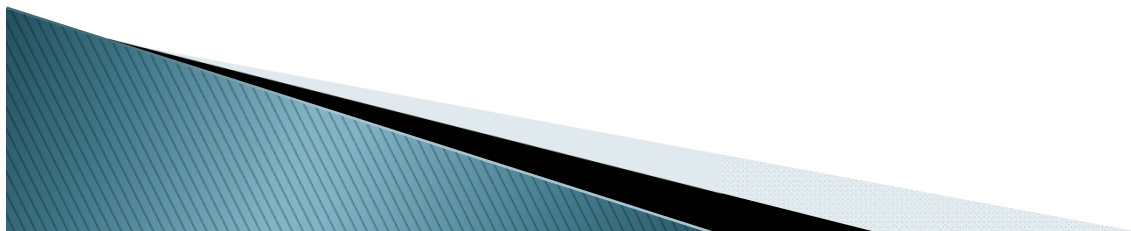
Second struct pre-allocated global memory, and function that assigns it to MpegEncContext at the point where MpegEncContext is initialized

```
MpegEncContext_Persistent_ pre_allocation[4] ;  
int MpegEncContext_Persistent_Allocated_Count = 0;  
  
void assign_allocated_Persistent_struct(void * s_struct)  
{  
    MpegEncContext * ptr = (MpegEncContext *) s_struct;  
    ptr->Persistent_ = & testGlobal[MpegEncContext_Persistent_Allocated_Count ++ ] ;  
}
```

Extend Analysis to all heavily used structs in the FFmpeg

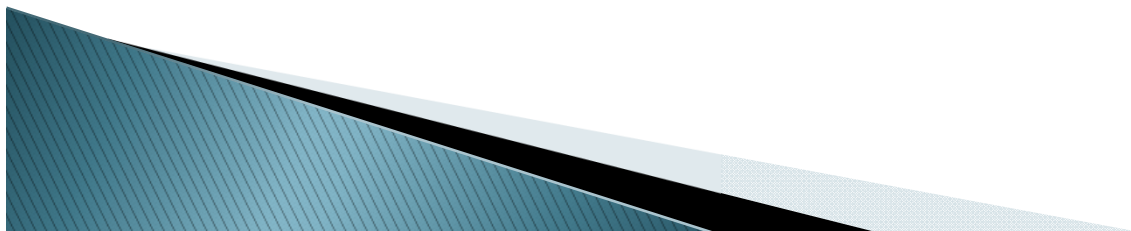
▶ Challenges

- The FFmpeg had over 550,000 lines of C codes.
- Each struct had a lot of members, which had to be traced individually, eg. The MpegEncContext struct had about 200 member to be traced.
- Some structs had members with same name, for this reason a simple find and replace may not work.



Solution – Created a custom code edition App

- ▶ Searching with complex criteria.
- ▶ Execute build command, and with one click opens up the source file and scrolls to error line.
- ▶ Automatically place variable access counters at the appropriate position in the source code



SDM Code Editor
Load Files | Batch Jobs | Simple Search | Run

Search Criteria

AND OR

AND OR

AND OR

Search

Display File Content Selection

Manual Code Editing

Read

Write

replace add to end of line Use Keyboard shortcut

```

1374 /* redraw edges for the frame if decoding didn't complete */
1375 // just to make sure that all data is rendered.
1376 if (CONFIG_MPEG_XVMC_DECODER && s->avctx->xvmc_accel
1377     ff_xvmc_field_end(s);
1378 } else if ((s->error_count || s->encoding || !(s->avctx->codec->capa
1379           !s->avctx->hwaccel &&
1380           !(s->avctx->codec->capabilities & CODEC_CAP_HWACCE
1381           s->unrestricted_mv &&
1382           s->current_picture.f.reference &&
1383           !s->intra_only &&|
1384           !(s->flags & CODEC_FLAG_EMU_EDGE)) {
1385     int hshift = av_pix_fmt_descriptors[s->avctx->pix_fmt].log2_chrc
1386     int vshift = av_pix_fmt_descriptors[s->avctx->pix_fmt].log2_chrc
1387     s->dsp.draw_edges(s->current_picture.f.data[0], s->current_pi
1388         s->Persistent_->h_edge_pos, s->Persistent_->v_edg
1389         EDGE_WIDTH, EDGE_WIDTH,
1390         EDGE_TOP | EDGE_BOTTOM);
1391     s->dsp.draw_edges(s->current_picture.f.data[1], s->current_pi
1392         s->Persistent_->h_edge_pos >> hshift, s->Persistent_
1393         EDGE_WIDTH >> hshift, EDGE_WIDTH >> vshift,
1394         EDGE_TOP | EDGE_BOTTOM);
1395     s->dsp.draw_edges(s->current_picture.f.data[2], s->current_pi
1396         s->Persistent_->h_edge_pos >> hshift, s->Persistent_

```

```

demo@ubuntu:~/share/ff$ make
CC libavcodec/mpegvideo.o
libavcodec/mpegvideo.c: In function âtomtom_log_mem_allocate_Persistentâ:
libavcodec/mpegvideo.c:149: warning: ISO C90 forbids mixed declarations and code
libavcodec/mpegvideo.c: In function âff_MPV_frame_endâ:
libavcodec/mpegvideo.c:1383: error: âMpegEncContextâ has no member named âintra_onlyâ
libavcodec/mpegvideo.c: In function âMPV_decode_mb_internalâ:
libavcodec/mpegvideo.c:2342: error: âMpegEncContextâ has no member named âintra_onlyâ
libavcodec/mpegvideo.c: In function âff_draw_horiz_bandâ:
libavcodec/mpegvideo.c:2571: error: âMpegEncContextâ has no member named âintra_onlyâ
make: *** [libavcodec/mpegvideo.o] Error 1
demo@ubuntu:~/share/ff$

```

Extract Error Lines

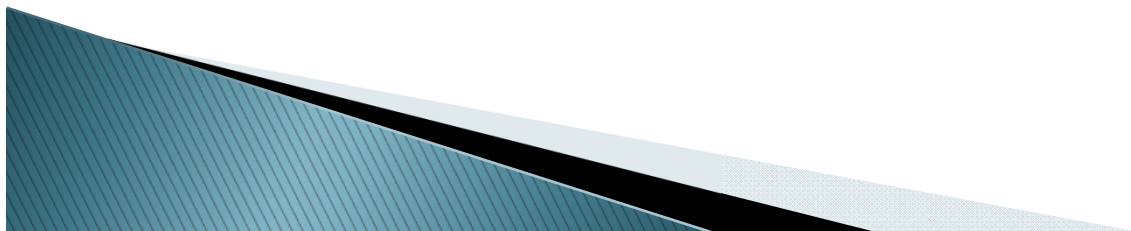
```

demo@ubuntu:~/share/ff$ make
CC libavcodec/mpegvideo.o
libavcodec/mpegvideo.c: In function âtomtom_log_mem_allocate_Persistentâ:
libavcodec/mpegvideo.c:149: warning: ISO C90 forbids mixed declarations and code
libavcodec/mpegvideo.c: In function âff_MPV_frame_endâ:
libavcodec/mpegvideo.c:1383: error: âMpegEncContextâ has no member named âintra_onlyâ
libavcodec/mpegvideo.c: In function âMPV_decode_mb_internalâ:
libavcodec/mpegvideo.c:2342: error: âMpegEncContextâ has no member named âintra_onlyâ
libavcodec/mpegvideo.c: In function âff_draw_horiz_bandâ:
libavcodec/mpegvideo.c:2571: error: âMpegEncContextâ has no member named âintra_onlyâ
make: *** [libavcodec/mpegvideo.o] Error 1

```


Conclusion

- ▶ Application optimization on hybrid memory
- ▶ Supporting tools
 - DASP : dataflow tool
 - Compiler approach
 - Interactive approach



Thank you!

